Adaptive Performance Optimization for Large-Scale Traffic Control Systems

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Abstract—In this paper, we study the problem of optimizing (fine-tuning) the design parameters of large-scale traffic control systems that are composed of distinct and mutually interacting modules. This problem usually requires a considerable amount of human effort and time to devote to the successful deployment and operation of traffic control systems due to the lack of an automated well-established systematic approach. We investigate the adaptive fine-tuning algorithm for determining the set of design parameters of two distinct mutually interacting modules of the traffic-responsive urban control (TUC) strategy, i.e., split and cycle, for the large-scale urban road network of the city of Chania, Greece. Simulation results are presented, demonstrating that the network performance in terms of the daily mean speed, which is attained by the proposed adaptive optimization methodology, is significantly better than the original TUC system in the case in which the aforementioned design parameters are manually fine-tuned to virtual perfection by the system operators.

Index Terms—Adaptive control, large-scale systems, parameter optimization, self-optimizing control, traffic control.

I. INTRODUCTION

Currently, a considerable amount of human effort and time is spent on the calibration of operations of large-scale traffic control systems (LSTCSs). Minor changes in the transport system infrastructure, e.g., installing a new variable message sign in a motorway network, modifying the traffic light signal phasing at an urban junction, and deploying a new bus in a public transport system or a new automated guided vehicle (AGV) in a seaport container terminal, may require the involvement of significant human effort and time to readjust and reprogram the LSTCS decision-making mechanisms. Moreover, the continuous medium- and long-term variations of the overall transport system dynamics (e.g., due to changes of traffic demand or the number of passengers who use the particular transport system) call for a frequent or even continuous maintenance of LSTCSs, which, if properly done, makes the maintenance of LSTCSs extremely costly. In several cases, the result is that system maintenance is neglected and the system performance deteriorates year after year.

In general, the same process that is required for calibration may also be used in the initial fine-tuning of the control system during its first installation. Both tasks (initial fine-tuning and calibration) are performed (if at all) by experienced personnel due to the lack of an automated and systematic approach; thus, there is no guarantee that the overall fine-tuning or maintenance will successfully end up. In some cases, the LSTCS has never achieved a satisfactory performance in the first place, e.g., in the reported case of the traffic signal control strategy split cycle and offset optimization technique (SCOOT), which is the most popular traffic management system worldwide, in the city of Nijmegen, The Netherlands [1], where the SCOOT application was completely abandoned in the end.

Urban and motorway traffic control systems, LSTCSs for public transport systems, and LSTCSs for railway, airport, and seaport operations are all specific examples of LSTCSs that call for calibration while the system is in operation. In all of these systems, the maintenance involves the recalibration, readjustment, and reprogramming of hundreds of parameters, rules, operational schedules, and decision-making mechanisms, which influence the transport system operations in a highly complex manner. Moreover, while the system is in operation, the use of heuristics, trial-and-error methods, and experience-based techniques involves the risk of poor system performance over a lengthy period of time, which may lead to delays, severe congestion, poor quality of service (QoS), and increased greenhouse gas (GHG) emissions during this period. Finally, note that the involvement of the human factor for the installation, maintenance, and renovation of LSTCSs also involves the risk of unsafe operations: human mistakes due to the lack of expertise and exhaustive working conditions may lead to decisions or actions that put safety at stake.

Recently, we have introduced and analyzed a new family of algorithms, hereafter called adaptive fine-tuning (AFT), that can be used toward the development of a generic, efficient, and systematic approach for the automated fine-tuning of LSTCSs [2]–[5]. The main attributes of these algorithms may be summarized as follows.

- The algorithms are based on adaptive optimization (AO) principles, and as a result, they do not require any a priori knowledge or assumption on the traffic system dynamics; moreover, they can be implemented to any type of LSTCS,
In the last decades, attempts have been made, particularly LSTCS applications, to develop model-based, i.e., either theoretical- or simulation-based, designs that produce “good” sets of tunable parameters, and although these approaches have helped, in some cases, to reduce the time and effort for installation and maintenance, they did not manage to eliminate or at least significantly reduce the involvement of the human factor. One example in this class is the implementation of a variable-speed-limit system on the U.K. motorway M42 [9]. Although the initial tunable parameters of the system (which correspond to speed and flow activation/deactivation thresholds) were “optimized” using theoretical tools from the traffic flow theory and extensive simulation experiments, it took more than one year of the calibration of the aforementioned thresholds until the system has reached an acceptable performance, and during this initial deployment phase, the system performance was sometimes worse than the no-control case.

There have also been some attempts to incorporate optimization-based tools into the maintenance procedure; for example, see [10]–[15] and the references therein for an indicative list of references. In these cases, the problem of providing efficient maintenance is formulated as an optimization problem, where the tunable LSTCS parameters are chosen to optimize a performance criterion (e.g., average network speeds in traffic networks, average delays in airborne or seaborne transport systems, the total number of containers that are loaded or unloaded in seaport container terminals, and the average deviation from the operational schedules in public transport systems). However, the optimization of such a performance criterion requires perfect or, at least, very accurate knowledge of the transport system dynamics and the demand. To deal with this problem, optimization-based approaches employ simulation-based or theoretical models to represent the actual system dynamics. Then, assuming that these models quite accurately represent the actual LSTCS operations, different optimization algorithms (e.g., gradient-descent, Gauss–Newton, evolutionary programming, or neural-network-based optimization algorithms) are applied to extract the optimal values of the tunable parameters. However, these approaches have the following two disadvantages: 1) They require extensive and continuous calibration of the simulation/theoretical-based models to optimize their approximation accuracy with respect to the actual transport system operations, and 2) they face the tradeoff between simplicity and accuracy, and in most cases, accuracy has to be sacrificed to avoid the use of extreme computational requirements of simulation or mathematical models that employ a detailed modeling of the LSTCS operations.

B. Adaptive and Neural/Learning Methods

One possible way of bypassing the aforementioned problems is to incorporate adaptive or adaptive-like designs (e.g., neural, fuzzy, or iterative learning methods) to update the parameters of the LSTCS, which render several advantages contrary to the simulation/theoretical-based techniques. AFT belongs to the family of the so-called AO methods, e.g., the simultaneous perturbation stochastic approximation (SPSA) [16], [17]. These methods probably provide the most promising approach for the development of a systematic methodology for automatic, safe, robust, and efficient maintenance and renovation of LSTCSs. The basic functioning procedure for AO methods may be summarized as follows (see Fig. 1).

- At the end of appropriately defined periods (e.g., at the end of each day), the AO algorithm receives the value of...
Unfortunately, these designs suffer from the following two severe drawbacks.

1) Although there are some AO methods that utilize a neural network training phase to capture past knowledge (for example, see [19]), the majority of AO methods do not have any mechanism for incorporating the knowledge captured in the past with regard to the dependence of the LSTCS performance on the tunable parameters and the external factors (demand). In case that such dependence is highly nonlinear and complex, the aforementioned algorithms fail to produce any improvement of the overall LSTCS performance when applied to the baseline control modules under study (see [2] for details).

2) Most importantly, the use of random perturbations in the AO algorithms may lead to an unacceptable value of the LSTCS performance; even a small perturbation of a “good” set of tunable parameters may lead to an unacceptable or, even worse, an unstable or catastrophic behavior. Hence, AO methods possess the disadvantage of not guaranteeing efficient and, most importantly, safe performance at the perturbation phase (in the method described in [19], offline simulation experiments are used to predefine acceptable levels of parameter perturbations for a stable traffic behavior).

In a series of papers [2]–[5], we introduced and analyzed a new family of AO algorithms that can overcome the aforementioned limitations (1) and (2). This approach appropriately combines the favorable features of AO algorithms with the features of approximation and adaptive mechanisms to come up with an AO methodology that can rapidly and efficiently optimize systems of arbitrary complexity and scale such as LSTCS and, most importantly, guarantee robust and safe performance while the maintenance operation is on.

### III. Problem Formulation

Consider a general discrete-time control system where the underlying dynamics are described according to the following nonlinear first-order difference equation:

\[
z_{t+1} = F(z_t, u_{i,t}, d_t, t), \quad z_0 = z(0)
\]

where \(z_t, u_{i,t}, \) and \(d_t\) are the vectors of system states, control inputs, and exogenous (possibly measurable) signals, respectively, \(t\) denotes the discrete-time index, \(i\) denotes the controller index, and \(F(\cdot)\) is a sufficiently smooth nonlinear vector function. Note that the proposed methodology can be applied to a system, although the function \(F\) is unknown. In addition, consider that one or more control laws are applied to the system (1), which are described as follows:

\[
u_{i,t} = \varphi_i(\theta_i, z_t)
\]

where \(\varphi_i(\cdot)\) are known smooth vector functions, and \(\theta_i\) is the vector of the \(i\)th controller tunable parameters. Note that we do not impose any restriction on either the form of (2) or the number of the applied control laws. In addition, the discrete-time index \(t\) may be different for each control law \(i\).
The overall system performance is evaluated through the following objective function (performance index):

\[
J(\theta; z_0, D_T) = \pi_T(z_T) + \sum_{i=1}^{I} \pi_{i,t}(z_i, u_{i,t})
\]

\[
= \pi_T(z_T) + \sum_{i=1}^{I} \sum_{t=0}^{T-1} \pi_{i,t}(z_i, \varnothing(\theta_i, z_i))
\]

where \( \theta = \text{vec}(\theta_1, \theta_2, \ldots, \theta_I), \pi_T, \) and \( \pi_{i,t} \) are known nonnegative functions, \( J \) is the number of the fine-tuned controllers, \( T \) is the finite-time horizon over which the control laws (2) are applied, and \( D_T = [d_0, d_1, \ldots, d_{T-1}] \) denotes the time history of the exogenous signals over the optimization horizon \( T \). By defining \( x = \text{vec}(z_0, D_T) \), (3) may be rewritten as

\[
J(\theta; z_0, D_K) = J(\theta, x).
\]

Equation (4) indicates that the system performance is affected by the vector of the tunable parameters \( \theta \). The problem at hand is to develop an appropriate iterative algorithm, which will be applied every \( T \) and will update the current control system parameters vector \( \theta \) to achieve better performance but also provide a safe and efficient behavior. This condition means that the algorithm should guarantee a stable sustainable system performance.

In every iteration \( k \) of the algorithm (fine-tuning experiment), the following two cases take place.

- The current controller design parameter vector \( \theta_k \) through the measurement

\[
J_k \equiv J(\theta_k, x_k).
\]

- The current controller design parameter vector \( \theta_k \) is updated so that it converges, as close as possible, to one of the local minima \( \theta^* \) of the average value of \( J \) (with respect to the exogenous random vectors \( x_k \)), which is defined according to

\[
\frac{\partial}{\partial \theta} E [J(\theta^*, x_k)|G_k] = 0
\]

where \( G_k \) is an appropriately defined term that refers to the past values of vector \( \theta \) and the exogenous inputs \( x \).

The requirement of the convergence of \( \theta_k \) to one of the local minima \( \theta^* \) is not sufficient in most practical situations. In addition to this requirement, the fine-tuning algorithm should provide safe and efficient performance during the fine-tuning. More precisely, at each iteration of the fine-tuning algorithm, the performance index measurement should satisfy

\[
J_k \leq J_{k-1} + \epsilon_k
\]

where \( \epsilon_k \) is an appropriately defined positive term, whose magnitude is proportional to the magnitude and variance of the exogenous inputs.

The requirement (7) is more than crucial in most practical LSTCS fine-tuning applications, because the violation of such a requirement may cause serious performance and safety problems. For example, in the case of the fine-tuning of traffic control systems, the violation of requirement (7) may lead to serious problems (e.g., complaints and dangerous driving), which may force the traffic operators to cancel the fine-tuning. Similarly, in the case of the fine-tuning of LSTCSs for mechanical structures, the violation of this requirement may cause the permanent deformation or even the destruction of the structure. Note that standard AO methodologies, e.g., the SPSA algorithm, cannot guarantee that the requirement (7) holds during the fine-tuning mainly due to the use of random perturbations of the controller parameters.

IV. ADAPTIVE FINE-TUNING ALGORITHM

This section briefly presents the main structure and the performance characteristics of the AFT algorithm.

A. Structure of the AFT Algorithm

The main components of the employed algorithm are summarized as follows.

- An approximator \( \hat{J}(\theta, x) \) is used (e.g., a neural network or a polynomial-like approximator) to obtain an approximation of the nonlinear mapping \( \hat{J}(\theta, x) = J(\theta, x) \).

- An online adaptive/learning mechanism is employed to “train” the aforementioned approximator. Globally convergent learning algorithms (for example, see [25] and [26]) are required for such a purpose.

- At each algorithm iteration \( k \), several randomly chosen candidate perturbations of vector \( \theta_k \) are selected, and the effect of each of these perturbations to the LSTCS performance is estimated using the aforementioned approximator. The perturbation that corresponds to the “best” estimate (i.e., the perturbation that leads to the best value for \( J \)) is picked to be the new tunable parameter values \( \theta_{k+1} \), which will be applied at the next period (e.g., the next day).

B. Performance Approximator

As aforementioned, for the approximation of the objective function \( J(\theta, x) \), a polynomial-like approximator with \( L_g \) regressor terms is used, which takes the form

\[
\hat{J}(\theta, x) = \varnothing^T \phi(\theta, x)
\]

where \( \varnothing \) denotes the matrix of the approximator parameter estimates, and

\[
\phi(\theta, x) = [\phi_1(\theta, x), \phi_2(\theta, x), \ldots, \phi_{L_g}(\theta, x)]^T.
\]

The nonlinear functions \( \phi_i(\theta, x) \) are given by

\[
\phi_i(\theta, x) = \Sigma d_i(\theta_{m_i}) \cdot \Sigma d_i(\theta_{m_2}) \cdot \Sigma d_i(\theta_{m_3}), \quad d_i \in \{0, 1\}
\]

where \( d_i \), \( m_i \) are randomly chosen at each iteration of the AFT algorithm, and \( \Sigma(\cdot) \), \( \Sigma(\cdot) \) are smooth monotone nonlinear functions. In the neural-networks literature [27], [28], these
functions are usually chosen to be “sigmoidal.” In this application, we choose
\[ S(\theta) = \tanh(\lambda_1 \theta + \lambda_2), \quad \tilde{S}(x) = \tanh(\lambda_3 x + \lambda_4) \]  
(11)
where \( \lambda_i \) are nonnegative real numbers that are initially defined by the user. After four or five iterations of the algorithm, the values of \( \lambda_i \) are optimized to minimize \( \min_{\ell=1}^{k-1} (J_\ell - \bar{\theta}_\ell^T \phi_\ell) \).

C. Algorithm Design and Convergence Properties

Table I presents a description of the design parameters and variables used within the AFT algorithm, whereas Table II presents a mathematical description of the AFT dynamics.

The proposed algorithm assumes that an estimate or prediction \( \bar{x}_k \) of the vector \( x_k \) is available. In several applications such an assumption is realistic, because the entries of \( x_k \) correspond to system states and exogenous inputs that are available for measurement. However, there may be cases where such an assumption is not realistic. In this case, \( \bar{x}_k \) can be estimated or predicted using appropriate estimation algorithms.

Note that, similarly to conventional AO algorithms, the proposed algorithm introduces random perturbations to the current control design parameter vector \( \theta \). In addition, the use of random perturbations is crucial for the efficiency of the proposed algorithm, because it provides the so-called persistence of excitation (PE) property, which is a sufficient and necessary condition for the neural approximator \( \hat{J} \) to efficiently learn the unknown function \( J \). However, due to the use of step 6 (see Table II), the proposed methodology avoids poor performance or instability problems and guarantees safe and efficient performance, because (7) is fulfilled.

Contrary to other applications of neural approximators, where the number of neurons \( L_g \) should be large enough to guarantee efficient approximation over the whole input set, this is not the case here. In the case of the proposed algorithm, it is sufficient that the approximator has enough regressor terms to come up with an approximation of the unknown function \( J \) over a “small neighborhood” around the most recent vector \( \theta_k \).

As shown in [2]–[5], using strict mathematical arguments, if the structure of the approximator and its learning mechanism satisfy certain design considerations (that are independent of the particular application), then the aforementioned process guarantees rapid convergence of the overall maintenance procedure to the same performance levels that would have been obtained if efficient nonlinear optimization schemes such as the steepest descent or Gauss–Newton schemes can be applied to the particular problem. Most importantly, the aforementioned procedure guarantees safe, stable, and efficient transient performance, because the system performance during maintenance remains within acceptable levels that can be, in the worst case, similar to the system performance before the maintenance has started.

V. TRAFFIC-RESPONSIVE URBAN CONTROL SIGNAL CONTROL STRATEGY

The TUC signal control strategy (see [6]–[8] for details) is a recently developed efficient TUC strategy whose design principles are based on the feedback control theory as opposed to most of the existing strategies that employ model-based optimization techniques. TUC consists of four distinct interconnected control modules that allow for real-time control of the following traffic measures:

1) green times (the split module);
2) the cycle time (the cycle module);
3) the offset (the green wave along an arterial);
4) the provision of public transport priority.
These four control modules are complemented by the fifth data-processing module. All control modules are based on feedback concepts of various types, which leads to TUC’s computational simplicity compared to model-based optimization techniques.
approaches, without sacrificing efficiency. In this paper, we will concentrate on the fine-tuning of the design parameters of the following two distinct mutually interacting control modules of TUC: 1) split and 2) cycle. Note that the proposed algorithm can also be applied to the fine-tuning of other control modules.

In the next two sections, we present the control laws that govern the split and cycle modules of TUC.

A. Split Control Module

The split control part of the TUC signal control strategy is derived from a formulation in the format of a linear–quadratic (LQ) control problem that leads to the multivariable regulator

$$g_t = g^N - Lz_{t_s}$$

where $t_s = 0, 1, 2, \ldots$ is the discrete time index, with the sample time period typically equal to the cycle time duration $C$, $g_t$ is the control vector (that consists of the green times of all stages in all junctions) that will be applied during the next cycle, $g^N$ is a nominal control vector (that consists of the nominal green times) that corresponds to a prespecified fixed signal plan (the impact of this plan on the resulting control was found to be limited), and $z_{t_s}$ is the state vector (that consists of the vehicle numbers in all network links during the last cycle) that is estimated by the data-processing module of TUC. Finally, $L$ is a constant feedback gain matrix (of appropriate dimensions) that is calculated offline based on a straightforward procedure according to the LQ regulator methodology. The entries of matrix $L$ depend on the network geometry, the turning rates, and the saturation flows. The sensitivity of TUC’s performance to moderate variations of these values has been found to be negligible [6], [29]. The aim of (12) is to balance the relative space occupancies $z_i/z_{i,\text{max}}$ in the network links to minimize the risk of queue spillovers, which may lead to a waste of green time and even to gridlocks. To this end, the regulator (12) may apply an inherent gating, i.e., reduce the green time of links that feed a saturating road, although these links are two or more junctions away.

The number of vehicles $z_{i,t_s}$ for link $i$ during the last cycle is estimated through the following equation:

$$z_{i,t_s} = z_{i,\text{max}}f(o_{i,t_s}, l_i)b_i$$

where $o_{i,t_s}$ denotes the measured average time occupancy (usually measured by loop detectors that are located at a certain distance from the stop line) during the last cycle time, $f(\cdot)$ is an empirical function [8], [30] that is constructed from practical investigations, and $l_i$ denotes the distance of the loop detector from the stop line divided by the total link length. Finally, $b_i$ is a nonnegative design parameter for each link $i$ (the so-called “importance factor”), which is introduced such that the $z_{i,t_s}$ values that result from (13) are multiplied with the corresponding $b_i$ before being used in the multivariable regulator (12). The default values are $b_i = 1$, but experienced system operators may manually select a real value $b_i \in (0, 3]$ to increase or decrease the importance of specific links, i.e., make them look more or less saturated than the measurements actually reflect. These design parameters are critical for the successful deployment and operation of the TUC signal control strategy and were hence selected for automated fine-tuning by the AFT algorithm.

B. Cycle Control Module

Cycle control is another module that TUC uses to influence traffic conditions. Longer cycle times typically increase the capacity of a junction but, on the other hand, may increase vehicle delays in undersaturated junctions due to longer waiting times during the red phase or, even worse, create queue spillovers. Considering the aforementioned remarks, the objective of the cycle control module is to increase the junctions’ capacities as much as necessary to limit the maximum observed saturation level in the network. Within TUC, this objective is effectuated through the application of a simple feedback-based regulator that uses the current saturation level of a prespecified percentage of the network links as a criterion for the increase or decrease of the cycle. The cycle module control law takes the form

$$C_{t_c} = \begin{cases} C^N + K_1(\sigma_{t_c} - \sigma_{N_1}), & \text{if} \; \sigma_{t_c} \leq \sigma_{cr} \\ C^N - K_2(\sigma_{t_c} - \sigma_{N_2}), & \text{if} \; \sigma_{t_c} > \sigma_{cr} \end{cases}$$

where $t_c = 0, 1, 2, \ldots$ is the discrete-time index of the cycle control, $C^N$ denotes a nominal network cycle time (typically equal to the minimum permissible cycle $C_{\text{min}}$); $\sigma_{t_c}$ is a vector that is composed of the mean values of the space occupancies for the prespecified links over the last cycle control period; $\sigma_{N_1}$, $\sigma_{N_2}$, and $\sigma_{cr} \in [0, 1]$ denote user-defined design parameters; and $K_1$ and $K_2 > 0$ are networkwide design parameters, the selection of which affects the intensity of the cycle control module reactions and may hence cause the degradation of the overall performance of the TUC strategy if not suitably configured. In other words, high $K_1$ and $K_2$ values force the control law to strongly react even for small differences of $\sigma_{t_c}$ from $\sigma_{N_i}$, $i = 1, 2$. For this reason, the design parameters $K_1$, $K_2$, $\sigma_{N_1}$, $\sigma_{N_2}$, and $\sigma_{cr}$ were selected for automated fine-tuning by the AFT algorithm. After the application of (14), the calculated cycle time $C_{t_c}$ is constrained within the range $[C_{\text{min}}, C_{\text{max}}]$, if necessary, to become feasible, where $C_{\text{min}}$ and $C_{\text{max}}$ are the minimum and maximum permissible network cycle times, respectively.

VI. APPLICATION OF THE ADAPTIVE FINE-TUNING ALGORITHM TO THE TRAFFIC-RESPONSIVE URBAN CONTROL SIGNAL CONTROL STRATEGY

To evaluate the efficiency of the aforementioned AFT algorithm to the problem of optimizing the design parameters of the split and cycle control modules of TUC, extensive simulation experiments have been conducted. The performance of AFT is compared to the base case (the non-AFT case), where the aforementioned design parameters were manually fine-tuned to virtual perfection by the system operators for the original TUC system [30].

A. Network and Simulation Setup

For the simulation experiments of the proposed approach, the road network of the city center of Chania, Greece, was
considered. The model of the network (see Fig. 2) consists of 16 signalized junctions (nodes) and 60 links (arrows). Each network link corresponds to a particular junction phase. Typical loop-detector locations within the Chania urban network links are either around the middle of the link or some 40 m upstream the stop line. Note that severe congestion problems occur in the actual Chania network, which sometimes leads to grid-lock situations.

The commercial microscopic simulator AIMSUN (version 6.0.1) [31] was employed as a simulation tool. The simulation step for the microscopic simulation model of the urban road network of the city center of Chania, Greece, was set to 0.5 s. The traffic network characteristics (e.g., saturation flows and turning rates) and the fixed plan $g^N$ in (12) used in AIMSUN and TUC were suggested by the system operators of the traffic control center (TCC) of the city (details are omitted due to space constraints). Note that the fixed plan $g^N$ is one of the six fixed predefined network signal plans used by the TCC. For the application of the TUC strategy, the following typical design values were used.

- $t_s = C$.
- $t_c = 600$ s.
- $C_{\text{min}} = 60$ s.
- $C_{\text{max}} = 120$ s.
- $C^N = C_{\text{min}}$.

In addition, for the implementation of the AFT algorithm, the following design values were used.

- $T_h = 90$.
- $L_g = 150$.
- $K = 20$.
- $\alpha_s = \alpha = 0.1$.

Initial values were set to $\lambda_i$ according to $\lambda_1 = 100$, $\lambda_3 = 0.1$, and $\lambda_2 = \lambda_4 = 0$. In reference to the notation used in Section III, we have $\theta = \text{vec}(\theta_1, \theta_2)$, where $\theta_1 = (b_1, b_2, \ldots, b_{60})$ and $\theta_2 = (K_1, K_2, \sigma_{N_1}, \sigma_{N_2}, \sigma_{cr})$ are the design parameters of the split and cycle control modules of TUC, respectively. The initial values for these parameters were chosen to correspond to values that are usually chosen during the initial field implementation of the TUC system. More precisely, the parameters $\theta$ were initialized according to $\theta_1 = 1$ and $\theta_2 = (240, 300, 0.15, 0.6, 0.4)$. Finally, to assess the overall system performance, the criterion $J \equiv ms$ [see (3)] was set to the actual daily network mean speed.

To investigate the performance of the AFT algorithm under different traffic conditions, two basic traffic demand scenarios (time history of vehicles that enter the network in the network origins during the day) were designed based on actual measurements. The simulation horizon of each scenario is 4 h. Scenario 1 comprises medium demand in all network origins, whereas scenario 2 comprises high demand, and the network faces serious congestion for some 2 h, with some link queues spilling back into upstream links. For simplicity, we assume that a demand scenario with a time horizon of 4 h corresponds to a day. Each day (iteration of the AFT algorithm), a randomly perturbed 5% width version of the basic demand scenarios is produced, and the assessment criterion is gathered from the AIMSUN simulator. Then, the design parameters of TUC strategy are updated by the AFT algorithm according to the calculated assessment criterion.

The overall closed-loop scheme consists of the following two main control loops: 1) the inner loop and 2) the outer loop. The inner loop is used by the TUC strategy to produce the traffic signal settings. More specifically, at each cycle $C$, AIMSUN delivers the (emulated) occupancy measurements at the locations where detectors are placed (as in real conditions). These measurements are used by the control modules of the TUC strategy to produce the traffic signal settings (split and cycle). These signal settings are then forwarded to the microsimulator for application through the application programming interface (API) programming module of AIMSUN. The outer loop is used by the AFT algorithm to update the design parameters of the TUC strategy. More specifically, at each day, AIMSUN delivers the mean speed for the whole urban road network. The mean speed is used by the AFT algorithm to produce the design parameters of the split and cycle control modules of the TUC strategy (the vector $\theta = \text{vec}(\theta_1, \theta_2)$). These design parameters are then forwarded to the TUC strategy for application, and so forth.

Fig. 2. Chania urban road network.
scenario 1 and day 10 for scenario 2; see Figs. 3 and 4). Thus, if the learning period is excluded (see the last two columns in Table III), the improvement increases to 23% and 40% for scenarios 1 and 2, respectively.

Table IV displays two more assessment criteria that are gathered from the microscopic simulator AIMSUN for both simulation scenarios. The average delay time (DT) per kilometer traveled (in seconds per kilometer) and the average number of stops (NS) per vehicle and per kilometer traveled are illustrated. The table also presents the improvement of these criteria due to the use of the AFT algorithm (20.86% and 16.21% for scenario 1 and 37.72% and 21.59% for scenario 2).

Figs. 5 and 6 display the importance factors of the network links according to the optimal solution of the AFT algorithm for scenarios 1 and 2, respectively. The green color indicates low importance \((b_i \leq 0.7)\), the black color indicates medium importance \((0.7 < b_i < 1.3)\), and the red color indicates high importance \((b_i \geq 1.3)\). It is shown that the AFT algorithm increases the weight of the importance factors for network links along the main entrance to and exit from the city center (junctions 16, 15, 14, 13, 12, 5, and 4), whereas for other links, which are not very crucial for the overall system performance, the corresponding weights are decreased. Although there are links with the same color for both scenarios, there are other links that are green in scenario 1 and red in scenario 2, and vice versa. The AFT algorithm converges to a local optimal solution that fine-tunes the importance factors and optimizes the traffic control system performance, which also depends on the special characteristics of each demand scenario.

The AFT algorithm is less than 10 s, which means that the implementation of the algorithm in a real-time large-scale application will be feasible, regardless of the type of the operating traffic control system. Finally, note that the AFT algorithm can also be utilized as an offline network optimization tool for calculating optimum sets of design parameters for LSTCSs of any type, because its system dynamics and controls (1) and (2) and related performance criterion (4) incorporate all the necessary network characteristics.

C. Detailed Results

In this section, we report on some selected results, focusing on the city’s main shopping district (see Fig. 2, junction 5).
With regard to the split control module of the TUC strategy, Fig. 11 compares the time evolution of the design parameters $b_{15}$ and $b_{18}$ under the use of the AFT algorithm, with some optimized values for these parameters. The optimized values come from a manual fine-tuning procedure, previously performed by human experts in a field evaluation of the TUC strategy in the Chania network [30]. The manual tuning of $b_{15}$ and $b_{18}$ led to the optimized weights 1.8 and 0.6, respectively. This link weighting is quite reasonable, because link 15 is a crucial link in the main arterial of the city center, contrary to link 18, which does not carry substantial traffic loads. The AFT algorithm starts from the initial weights $b_{15} = b_{18} = 1$ and, by iteratively optimizing their values, converges to weights that are shown to be close to the roughly optimized values. Fig. 12 displays the aforementioned design parameters for demand scenario 2. The weights again converge close to the optimized values, although they are slightly different from scenario 1 due to different traffic conditions. What is clear in both figures is that link 15 is more important than link 18 for the network mean speed. Note that this case holds for several other network links that...
are not shown here (see Figs. 5 and 6 for a general view). This case demonstrates that the proposed algorithm is a feasible and viable solution for the automated parameter fine-tuning of such systems.

In the following discussion, we illustrate the impact of the AFT algorithm to the cycle control module of the TUC strategy.

The AFT algorithm, by changing the design parameters \( \theta_2 = (K_1, K_2, \sigma_{N_1}, \sigma_{N_2}, \sigma_{cr}) \), changes the cycle control feedback regulator (14) and succeeds in notably improving the mean speed of the simulated network. Fig. 13 displays different values of \( \sigma_{tc} \), the initial cycle control regulator, and the fine-tuned cycle control regulator after the convergence of AFT algorithm for scenarios 1 and 2, respectively. Recall that \( \sigma_{tc} \in [0, 1] \) is the average space occupancy for some prespecified percentage of network links over the last cycle control period. Here, we depict the three different regulators for \( \sigma_{tc} \) from 0 to 0.7.

The initial cycle control regulator [see (14)] consists of one monotonically increasing function for undersaturated traffic conditions \((\sigma_{cr} < 0.4)\) and one monotonically decreasing function for saturated traffic conditions \((\sigma_{cr} > 0.4)\). The maximum cycle \( C_{max} = 120 \text{ s} \) is applied for the critical occupancy \( \sigma_{cr} = 0.4 \). After the convergence of the AFT algorithm, a
cycle regulator with the following three regimes is obtained for both simulated scenarios: 1) one regime with increasing cycle periods for undersaturated traffic conditions; 2) one regime with maximum cycle periods \( C_{\text{max}} = 120 \, \text{s} \) (for \( 0.25 \leq \sigma_{cr} \leq 0.4 \) and \( 0.32 \leq \sigma_{cr} \leq 0.4 \) for scenarios 1 and 2, respectively) for traffic conditions that require us to increase the capacity of the network; and 3) one regime with decreasing cycle periods for saturated traffic conditions (see Fig. 13). For the third regime, the network faces severe congestion problems due to queue spillovers and partial gridlocks that lead to a strong performance deterioration. Again, there are slight differences between the fine-tuned cycle control regulators of scenarios 1 and 2, which depend on the different traffic characteristics.

In general, the derived trapezoidal shape (see Fig. 13) of the cycle control regulators over the saturation level of the network outperforms the initial cycle control regulator, as shown by the overall system performance. Eventually, the following three traffic regimes may be identified: 1) undersaturated traffic conditions; 2) critical traffic conditions; and 3) saturated traffic conditions. This way, the cycle control module of TUC applies appropriate cycle times for each case, i.e., smaller cycle times for regimes 1 and 3 and the maximum cycle time for regime 2, to maximize the network’s capacity. Note that regime 2 occurs for different values of \( \sigma_{cr} \) for demand scenarios 1 and 2, which means that the real-time implementation of the AFT algorithm can be vital for the overall system performance.

VII. CONCLUSION

This paper has investigated the efficiency of the AFT algorithm for the problem of optimizing the design parameters of LSTCSs that are composed of distinct mutually interacting modules. This AO methodology aims at replacing the conventional manually based optimization with a fully automated procedure. Extensive simulation experiments have been conducted for the signal control problem of the large-scale traffic network of the city of Chania, Greece, where the design parameters of two distinct mutually interacting modules of the TUC strategy were fine-tuned by the AFT algorithm. The simulation results and the comparison to the base case, where the aforementioned design parameters of the TUC system were manually fine-tuned to virtual perfection by the system operators, demonstrate the algorithm’s efficiency and feasibility.

The design parameters of the split and cycle control modules of TUC have been considered for fine-tuning. It was demonstrated that the application of the AFT algorithm to the TUC signal control strategy leads to better network performance (in terms of the daily mean speed) compared with the original TUC system. This case underlines the superiority of the fully automated optimization procedure, pursued by the AFT algorithm, even in the case that the design parameters are already manually fine-tuned by field experts.

With regard to the design parameters of the split control module, it was shown that the AFT algorithm increases the weight of the importance factors for network links along the main entrance to and exit from the city center, whereas for other links, which are not very crucial, the corresponding weights are decreased. Furthermore, it was shown that the AFT algorithm converges to a set of quite-reasonable design parameters that are close to the roughly optimized values provided by the system operators for the original TUC system. Finally, with regard to the design parameters of the cycle control module, it was shown that the AFT algorithm leads to cycle control regulators with a trapezoidal shape of three traffic regimes, which outperform the initial cycle control regulator.

Future work will deal with the application of the AFT algorithm to all control modules of the TUC strategy (i.e., split, cycle, and offset) and with investigations with regard to the following two approaches: 1) the online generation of the gain sequences \( \alpha_k \) and 2) the approximation of the objective function \( J(\theta, x) \) through support vector machine (SVM) regressors, which can possibly increase the efficiency of the algorithm.

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